

MOTIVATION

WANT:

- Cheap & representative data.
- Model charging for difficult subgroups.

REALITY:

- Data collection is expensive.
- Existing datasets are biased.

SOLUTION:

1. Combine existing data from data repositories.
2. Enforce group count requirements.

PROBLEM DEFINITION

GIVEN:

- Data sources with sampling costs.
- Groups with minimum count requirements.

GOAL: Minimize expected total query cost.

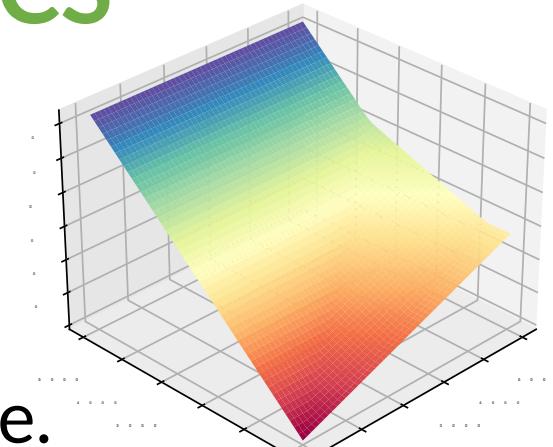
CONSTRAINT: Satisfy minimum count requirement.

QUERY MODEL: Uniformly random sampling.

KNOWN STATISTICS

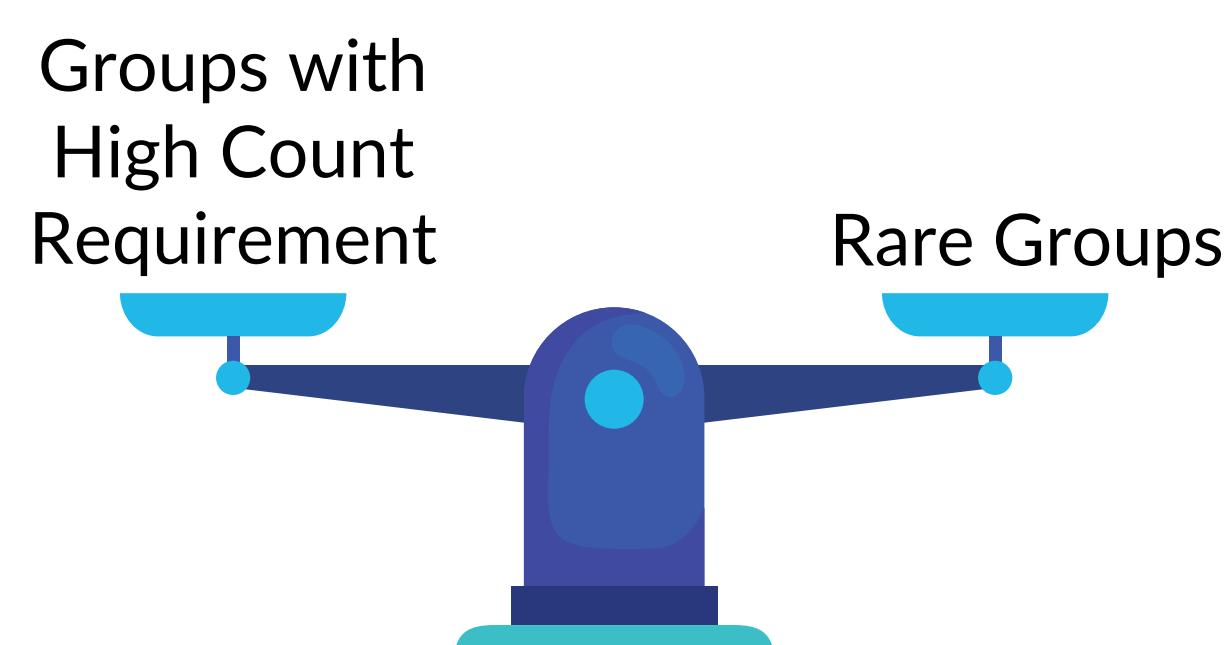
DYNAMIC PROGRAMMING

- Optimal but slow $O(Q^n)$ time.



HEURISTIC ALGORITHM RatioColl

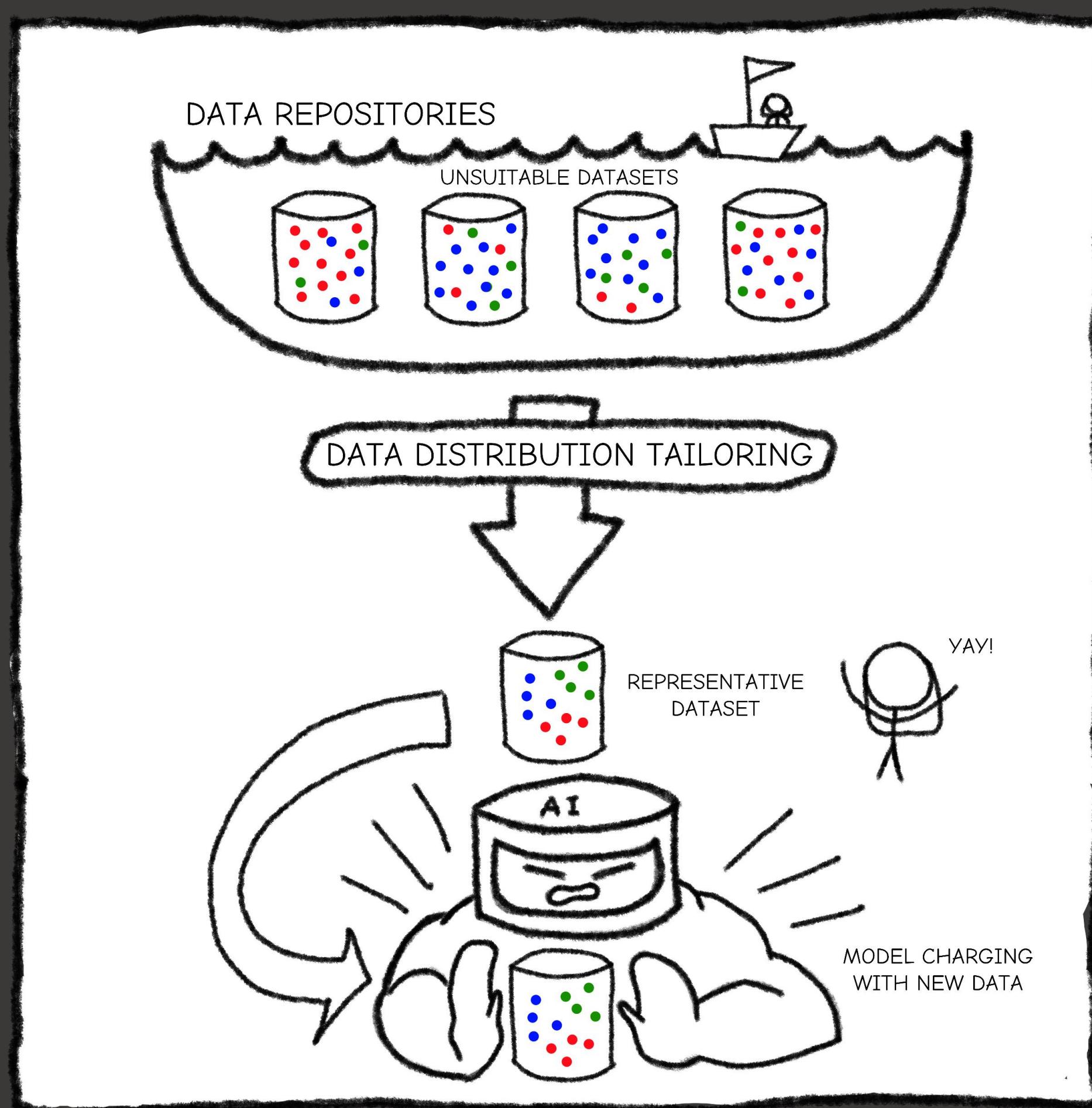
- Only \sqrt{Q} overhead cost in special case.
- Linear upper bound in general.



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Data Distribution Tailoring Revisited: Cost-Efficient Integration of Representative Data



UNKNOWN STATISTICS

ϵ -GREEDY BANDIT

- Same heuristic as RatioColl.
- No priors needed.
- Sublinear regret bound.



ALGORITHMS

RatioColl

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1:  $O \leftarrow \emptyset$ 
2: while query is not satisfied:
3:    $G^* = \operatorname{argmax}_{G_j, Q_j > 0} \left( Q_j \cdot \min_{i \in [n]} \left( \frac{c_i}{P_{i,j}} \right) \right)$  // Choose priority group
4:    $D^* = \operatorname{argmin}_{D_i} \left( \frac{c_i}{P_{i,*}} \right)$  // Maximize priority group per cost
5:    $s \leftarrow \operatorname{Query}(D^*)$ 
6:    $O \leftarrow O \cup \{s\}$ 
7: return  $O$ 

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EpsilonGreedy

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1:  $O \leftarrow \emptyset$ 
2: while query is not satisfied:
3:   if  $t \leq n$  then  $D^* \leftarrow D_t$  // Initialization
4:   else:
5:     with probability  $p = \sqrt[3]{\ln t / t}$ : // Exploration round
6:        $D^* \leftarrow$  random data source
7:     else: // Greedy exploitation round
8:       let  $\bar{P}_{i,j} \leftarrow$  estimate of  $P_{i,j}$  based on sample mean
9:        $R(G_j) \leftarrow \left( Q_j \cdot \min_{i \in [n]} \left( \frac{c_i}{\bar{P}_{i,j}} \right) \right)$  for each  $G_j$  // Reward of group
10:       $D^* \leftarrow \operatorname{argmax}_{D_i} \left( \frac{1}{C_i} \sum_{j \in [m]} \bar{P}_{i,j} \cdot R(G_j) \right)$  // Reward of data source
11:       $s \leftarrow \operatorname{Query}(D^*)$ 
12:       $O \leftarrow O \cup \{s\}$ 
13: update trackers for probability estimation
14: return  $O$ 

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SELECTED RESULTS

- RatioColl consistently out-performs SOTA.
- EpsilonGreedy competitive with SOTA despite needing no priors.

